

# Sorting

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Lecture 27

# Bubble sort method

Someone gives us a vector of unsorted numbers.

We want to obtain the vector sorted in ascending order.

- assign `IndexOfTheLastToCheck` be the *index* of the vector end
- 1 start sweeping from the beginning of the vector
- 2 Compare the 2 consequent elements till we reach the `IndexOfTheLastToCheck`
- 3 if the left element is larger we swap these 2 elements
- 4 move to the next pair to the right i.e. move to the **item 2**
  - notice that at the end of the sweep the *index* of the last element to check holds the largest element
  - so next sweep does not have to be that long.
  - it is shorter by one element
  - i.e. the *index* of the last element to check should be decreased by 1
- 5 decrease `IndexOfTheLastToCheck` by 1
- 6 if `IndexOfTheLastToCheck` > 1 repeat from the **item 1**

$x = [3, 1, 4, 5, 2]$

first sweep

$x = [\widehat{3}, 1, 4, 5, 2]$  swap

$x = [1, 3, 4, 5, 2]$  after swap

$x = [1, \widehat{3}, 4, 5, 2]$  no swap

$x = [1, 3, \widehat{4}, 5, 2]$  no swap

$x = [1, 3, 4, \widehat{5}, 2]$  swap

$x = [1, 3, 4, 2, \widehat{5}]$  sweep done

new sweep

$x = [1, \widehat{3}, 4, 2, 5]$  no swap

$x = [1, 3, \widehat{4}, 2, 5]$  no swap

$x = [1, 3, 4, \widehat{2}, 5]$  swap

$x = [1, 3, 2, 4, \widehat{5}]$  sweep done

new sweep

$x = [1, \widehat{3}, 2, 4, 5]$  no swap

$x = [1, 3, 2, \widehat{4}, 5]$  swap

$x = [1, 2, 3, 4, \widehat{5}]$  sweep done

last sweep

$x = [1, \widehat{2}, 3, 4, 5]$  no sweep

$x = [1, 2, 3, 4, \widehat{5}]$  sweep done



# Bubble sort properties

- This is the worst of all working algorithm!
- The execution time of this algorithm is  $\mathcal{O}(N^2)$
- Never use it in the real life!
- However it is very simple to program, and does not require extra memory for execution.

# Quick sort method

Much better yet simple algorithm

Let's discuss recursive realization

We will name our sorting function as `qsort`.

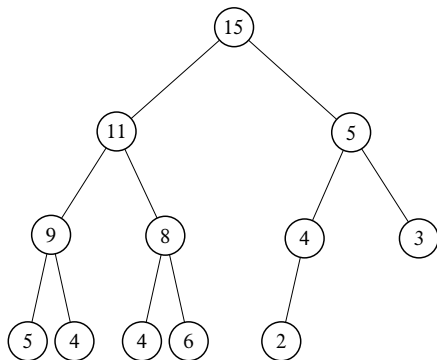
- choose a pivot point value
  - let's choose the pivot at the middle of the vector
  - $\text{pivotIndex} = \text{floor}(N/2)$
  - $\text{pivotValue} = x(\text{pivotIndex})$
- create two vectors which hold lesser and larger than `pivotValue` elements of the input vector.
- now concatenate the result of  
 $\text{xs} = [\text{qsort}(\text{lesser}), \text{pivotValue}, \text{qsort}(\text{larger})]$
- done

# Quick sort summary

- usually fast
- typical execution time  $\mathcal{O}(N \log_2 N)$
- but it is not guaranteed
  - However **for certain input vectors** execution time could be as long as  $\mathcal{O}(N^2)$

# Heap

Heap is a structure where parent element is larger or equal to its children.



The top most element of the heap is called root.

# Heap sorting method

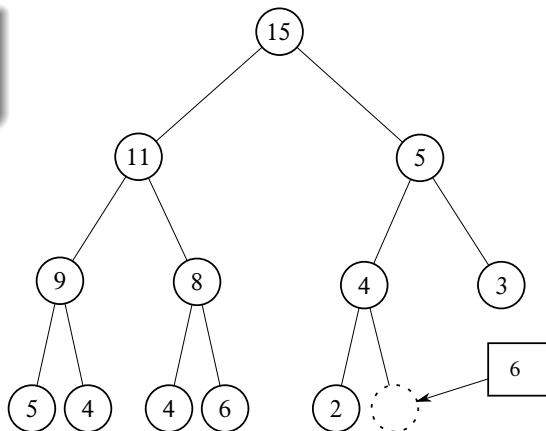
- 1 Fill the heap from the input vector elements
  - 1 take the element and place it at the bottom of the heap
  - 2 sift-up (bubble up) this element
  - 3 do the same with the next element
- 2 remove the root element since it is the largest
- 3 rearrange the heap i.e. sift-down
  - 1 take the last bottom element
  - 2 place it at the root
  - 3 check if parent is larger than children
    - 1 find the largest child element
    - 2 if the largest child is larger than parent swap them and repeat the check
- 4 repeat step 2 until no elements left in the heap

Heap sorting complexity  $\mathcal{O}(N \log_2 N)$

# Filling (sift-up) the heap

## Step 1

Place new element at the bottom of the heap

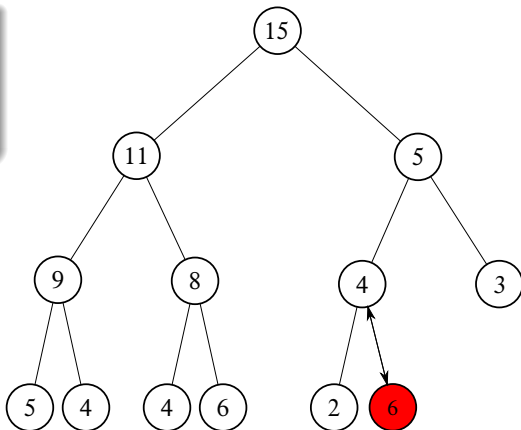




# Filling (sift-up) the heap

## Step 2

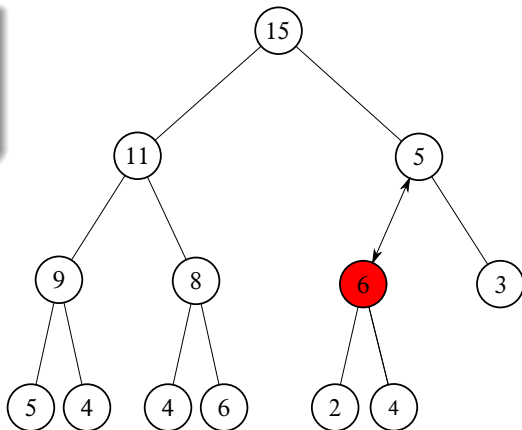
Check if parent is larger than child. If so swap them and repeat step 2.



# Filling (sift-up) the heap

## Step 2

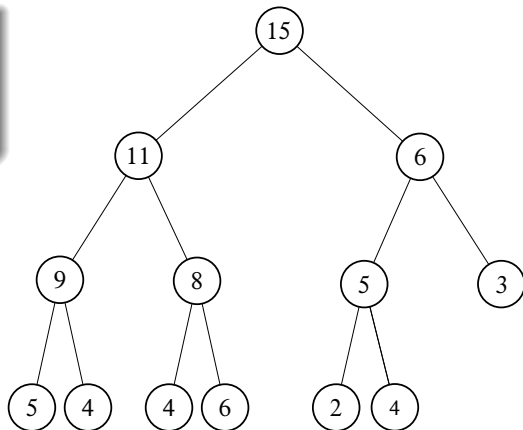
Check if parent is larger than child. If so swap them and repeat step 2.



# Filling (sift-up) the heap

## Step 2

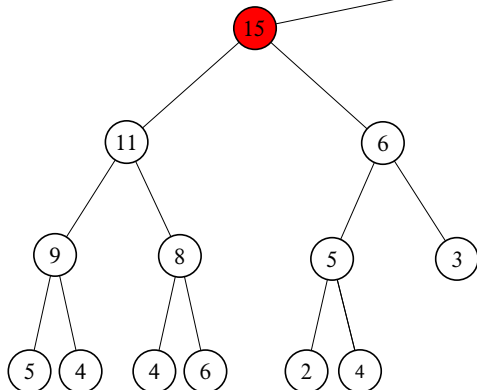
Check if parent is larger than child. If so swap them and repeat step 2.



# Removing from the heap (sift-down) the heap

Step 1

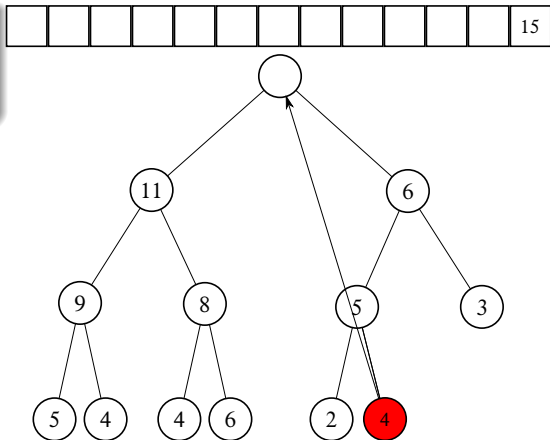
Remove the root element



# Removing from the heap (sift-down) the heap

## Step 2

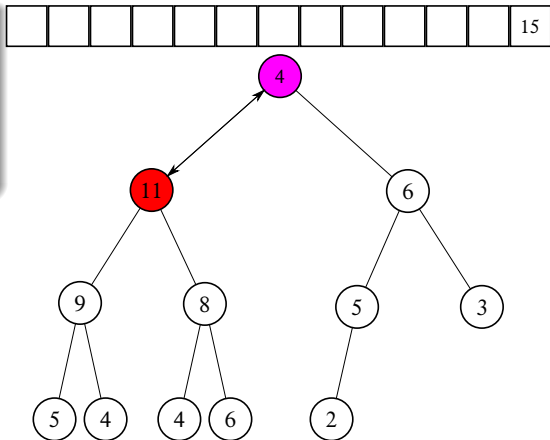
Place the last element of the heap to the root



# Removing from the heap (sift-down) the heap

## Step 3

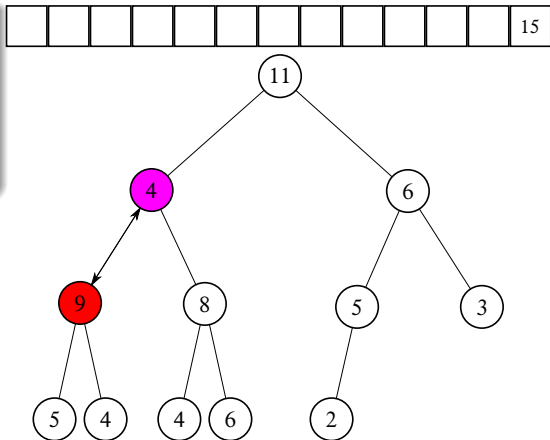
Check if parent is smaller than the largest child. If so swap and repeat step 3 else go to step 1



# Removing from the heap (sift-down) the heap

## Step 3

Check if parent is smaller than the largest child. If so swap and repeat step 3 else go to step 1

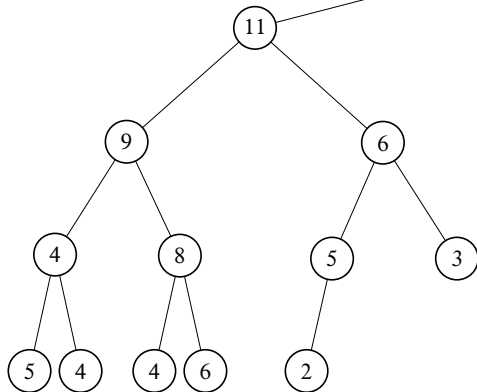
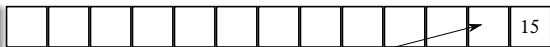


# Removing from the heap (sift-down) the heap

Sequence repeats

Step 1

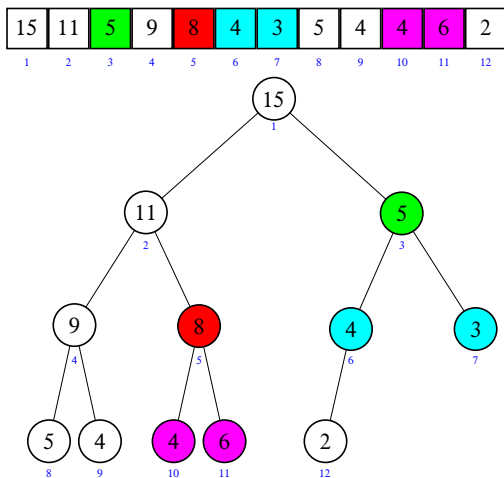
Remove the root element





# Vector heap representation

- Heap nodes are numbered consequently these numbers represent the node position in the vector.
- notice that parent and children have very simple relationship
  - if parent node index is  $i$ 
    - child 1 index is  $2i$
    - child 2 index is  $2i + 1$
  - if we know child index ( $i$ ) then
    - parent index is  $\text{floor}(i/2)$



# Matlab built in 'issorted'

Easy check if an array is sorted can be done with `issorted` which returns `true` or `false`.

```
>> x=[1,2,3];  
>> issorted(x)  
ans =  
1
```

`issorted` checks only for **ascending** order, for example

```
>> x=[3,2,1];  
>> issorted(x)  
ans =  
0
```

Recall that '0' is equivalent of `false` in Matlab