

# Sorting

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Lecture 07

# Bubble sort method

Someone gives us a vector of unsorted numbers.

We want to obtain the vector sorted in ascending order.

- assign `IndexOfTheLastToCheck` be the *index* of the vector end
- 1 start sweeping from the beginning of the vector
- 2 Compare the 2 consequent elements till we reach the `IndexOfTheLastToCheck`
- 3 if the left element is larger we swap these 2 elements
- 4 move to the next pair to the right i.e. move to the **item 2**
  - notice that at the end of the sweep the *index* of the last element to check holds the largest element
  - so next sweep does not have to be that long.
  - it is shorter by one element
  - i.e. the *index* of the last element to check should be decreased by 1
- 5 decrease `IndexOfTheLastToCheck` by 1
- 6 if `IndexOfTheLastToCheck` > 1 repeat from the **item 1**

$x = [3, 1, 4, 5, 2]$

first sweep

$x = [\widehat{3}, 1, 4, 5, 2]$  swap

$x = [1, 3, 4, 5, 2]$  after swap

$x = [1, \widehat{3}, 4, 5, 2]$  no swap

$x = [1, 3, \widehat{4}, 5, 2]$  no swap

$x = [1, 3, 4, \widehat{5}, 2]$  swap

$x = [1, 3, 4, 2, \widehat{5}]$  sweep done

new sweep

$x = [1, \widehat{3}, 4, 2, 5]$  no swap

$x = [1, 3, \widehat{4}, 2, 5]$  no swap

$x = [1, 3, 4, \widehat{2}, 5]$  swap

$x = [1, 3, 2, 4, \widehat{5}]$  sweep done

new sweep

$x = [1, \widehat{3}, 2, 4, 5]$  no swap

$x = [1, 3, 2, \widehat{4}, 5]$  swap

$x = [1, 2, 3, 4, \widehat{5}]$  sweep done

last sweep

$x = [1, 2, \widehat{3}, 4, 5]$  no sweep

$x = [1, 2, 3, 4, \widehat{5}]$  sweep done



# Bubble sort properties

- This is the worst of all working algorithm!
- The execution time of this algorithm is  $\mathcal{O}(N^2)$
- Never use it in the real life!
- However it is very simple to program, and does not require extra memory for execution.

# Quick sort method

Much better yet simple algorithm

Let's discuss recursive realization

We will name our sorting function as `qsort`.

- choose a pivot point value
  - let's choose the pivot at the middle of the vector
  - `pivotIndex=floor(N/2)`
  - `pivotValue=x(pivotIndex)`
- create two vectors which hold lesser and larger than `pivotValue` elements of the input vector.
- now concatenate the result of  
`xs=[qsort(lesser), pivotValue, qsort(larger)]`
- done

# Quick sort summary

- usually fast
- typical execution time  $\mathcal{O}(N \log_2 N)$
- but it is not guaranteed
  - However **for certain input vectors** execution time could be as long as  $\mathcal{O}(N^2)$