Root finding

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Lecture 06

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$$3x^3 + 2 = \sin x \rightarrow 3x^3 + 2 - \sin x = 0$$

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A general search algorithm is the following

- make a guess i.e. trial
- make intelligent new guess (x_{i+1}) judging from this trial (x_i)
- continue until $|f(x_{i+1})| > \varepsilon_f$ and $|x_{i+1} x_i| > \varepsilon_X$

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- some one think of any number between 1 and 100
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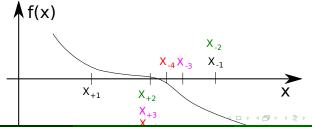
- some one think of any number between 1 and 100
- I will make a guess
- you provide me with either "less" or "more" depending where is my guess with respect to your number

How many guesses do I need? At most 7

Bisection method pseudo code

Works for any continuous function in vicinity of function root

- make initial bracket for search x_+ and x_- such that
 - $f(x_+) > 0$
 - $f(x_{-}) < 0$
- loop begins
- make new guess value $x_q = (x_+ + x_-)/2$
- if $|f(x_g)| \le \varepsilon_f$ or $|x_+ x_g| \le \varepsilon_X$ stop we found the solution with desired approximation
- otherwise if $f(x_g) > 0$ then $x_+ = x_g$ else $x_- = x_g$
- continue the loop



Bisection - simplified matlab implementation

```
function x sol=bisection(f, xn, xp, eps f, eps x)
% solving f(x)=0 with bisection method
 xq=(xp+xn)/2; % initial quess
 fg=f(xg); % initial function evaluation
 while ((abs(fq) > eps f) & (abs(xq-xp) > eps x))
   if (fq>0)
     xp=xq;
   else
    xn=xq;
   end
   xq=(xp+xn)/2; % update quess
   fg=f(xg); % update function evaluation
 end
 x sol=xq; % solution is ready
```

Bisection - example of use

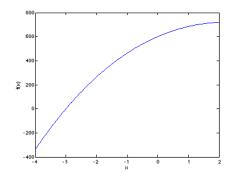
Let's define simple test function in the file 'function_to_solve.m'

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function ret=function_to_solve(x)
  ret=(x-10)*(x-20)*(x+3);
end
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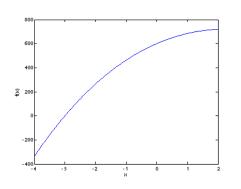
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```
bisection(...
@function_to_solve,...
-4.1, 2, ...
1e-6, 1e-8)
```

pay attention to the function handle operator @

$$ans = -3.0000$$

always cross check results

>> function_to_solve(ans)
ans = 3.0631e-07

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So what should we check for sure

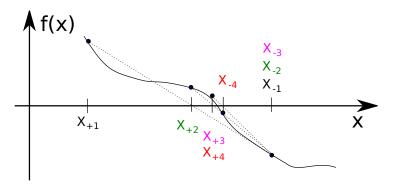
- f(xn) < 0
- ② f(xp) > 0

It would be handy to return secondary outputs

- with the value of function at the found solution point
- the number of iterations (good for performance tests)

False position (regula falsi) method

In this method we naively approximate our function as a line.



False position method - pseudo code

- make initial bracket for search x_+ and x_- such that
 - $f(x_+) > 0$
 - $f(x_{-}) < 0$
- loop begins
- draw a chord between points $(x_-, f(x_-))$ and $(x_+, f(x_+))$
- make new guess value at the point of the chord intersection with the 'x' axis

$$x_g = \frac{x_- f(x_+) - x_+ f(x_-)}{f(x_+) - f(x_-)}$$

- if $|f(x_g)| \le \varepsilon_f$ or $|x_+ x_g| \le \varepsilon_x$ stop we found the solution with desired approximation
- otherwise if $f(x_g) > 0$ then $x_+ = x_g$ else $x_- = x_g$
- continue the loop

Note: it looks like bisection except the way of updating x_g



We say that algorithm has defined convergence if it is possible to express

$$\lim_{k \to \infty} (x_{k+1} - x_0) = c(x_k - x_0)^m$$

Where x_0 is true root of the equation, c is some constant, and m is the order of convergence.

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Generally the speed of the algorithm is related to its convergence order. How ever other factors may affect the speed.